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APPLICATION NO.	F	ILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
10/082,776	02/25/2002		Victor A. Bennett	BENNETT 6-5 4410	
47396	7590	08/07/2006		EXAMINER	
HITT GAI	•		LI, AIMEE J		
AGERE SYSTEMS INC. PO BOX 832570 RICHARDSON, TX 75083				ART UNIT	PAPER NUMBER
				2183	
				DATE MAILED: 08/07/2006	

Please find below and/or attached an Office communication concerning this application or proceeding.

	Application No.	Applicant(s)				
	10/082,776	BENNETT ET AL.				
Office Action Summary	Examiner	Art Unit				
	Aimee J. Li	2183				
The MAILING DATE of this communication appears on the cover sheet with the correspondence address Period for Reply						
A SHORTENED STATUTORY PERIOD FOR REPLY WHICHEVER IS LONGER, FROM THE MAILING DOWN THE MAILING DOWN THE STATE OF THE MAILING DOWN THE STATE OF THE MAILING DOWN THE STATE OF THE MAILING DOWN THE MAILING THE MAILI	ATE OF THIS COMMUNICATION 36(a). In no event, however, may a reply be tin will apply and will expire SIX (6) MONTHS from a cause the application to become ABANDONE	N. nely filed the mailing date of this communication. (D) (35 U.S.C. § 133).				
Status						
	Responsive to communication(s) filed on 23 June 2006.					
· <u> </u>	,—					
•	Since this application is in condition for allowance except for formal matters, prosecution as to the merits is					
closed in accordance with the practice under Ex parte Quayle, 1935 C.D. 11, 453 O.G. 213.						
Disposition of Claims						
4) ⊠ Claim(s) <u>1-21</u> is/are pending in the application. 4a) Of the above claim(s) is/are withdray 5) □ Claim(s) is/are allowed. 6) ⊠ Claim(s) <u>1-21</u> is/are rejected. 7) □ Claim(s) is/are objected to. 8) □ Claim(s) are subject to restriction and/o	wn from consideration.					
Application Papers						
9) The specification is objected to by the Examine 10) The drawing(s) filed on is/are: a) accomposed and all accomposed and any objection to the Replacement drawing sheet(s) including the correct 11) The oath or declaration is objected to by the Example 11.	epted or b) objected to by the I drawing(s) be held in abeyance. Sec tion is required if the drawing(s) is ob	e 37 CFR 1.85(a). jected to. See 37 CFR 1.121(d).				
Priority under 35 U.S.C. § 119						
 12) Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f). a) All b) Some * c) None of: 1. Certified copies of the priority documents have been received. 2. Certified copies of the priority documents have been received in Application No. 3. Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)). * See the attached detailed Office action for a list of the certified copies not received. 						
Attachment(s)						
 Notice of References Cited (PTO-892) Notice of Draftsperson's Patent Drawing Review (PTO-948) Information Disclosure Statement(s) (PTO-1449 or PTO/SB/08) Paper No(s)/Mail Date 	4) Interview Summary Paper No(s)/Mail Do 5) Notice of Informal P 6) Other:					

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DETAILED ACTION

1. Claims 1-21 have been considered. Claims 1, 4, 8, 11, 15, and 18 have been amended as per Applicants' request.

Papers Submitted

2. It is hereby acknowledged that the following papers have been received and placed of record in the file: RCE as received 23 June 2006 and Amendment as received on 23 June 2006.

Claim Rejections - 35 USC § 103

- 3. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:
 - (a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the manner in which the invention was made.
- 4. Claims 1-3, 6-10, 13-14 are rejected under 35 U.S.C. 103(a) as being unpatentable over Parady, U.S. Patent Number 5,933,627 (herein referred to as Parady) in view of MicrowareTM's "OS-9: Real-Time Operating System The Complete Software Solution for Your Embedded Application" ©1999 (herein referred to as OS-9).
- 5. Referring to claim 1, Parady has taught a context switching system for a multi-thread execution pipeline loop having a pipeline latency, comprising:
 - a. A context switch requesting subsystem configured to:
 - Detect a device request from thread executing within said multi-thread execution pipeline loop for access to a device having a fulfillment latency exceeding said pipeline latency (Parady Abstract; column 1, lines 29-35

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and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3), and

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- ii. Generate a context Switch request for said thread (Parady Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3); and
- b. A context controller subsystem configured to receive said context switch request (Parady Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3).
- 6. Parady has not explicitly taught
 - a. A miss fulfillment first-in-first-out buffer (FIFO); and
 - b. Based thereon, store said thread in said miss fulfillment FIFO to prevent said thread from executing until said device request is fulfilled, said thread sequencing through said entire miss fulfillment FIFO before exiting therefrom.
- 7. However, Parady has taught using a round robin scheduling method (Parady column 4, lines 9-11). OS-9 has taught round robin scheduling (OS-9 page 5) with a
 - a. A miss fulfillment first-in-first-out buffer (FIFO) (OS-9 page 5); and
 - b. Based thereon, store said thread in said miss fulfillment FIFO to prevent said thread from executing until said device request is fulfilled, said thread sequencing through said entire miss fulfillment FIFO before exiting therefrom (OS-9 page 5).
- 8. In regards to OS-9, the event queue stores processes that are blocked, waiting for some event to complete, like an I/O access. This is similar to a miss fulfillment FIFO, which, by its name and the language in the claim, stores processes that wait for some other event, like an I/O

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access, to complete. A person of ordinary skill in the art at the time the invention was made would have recognized that a round-robin scheduling method is a fairly simple method that guarantees that all processes will be executed while ensuring that priority between processes is maintained, thereby preventing starvation of a process while giving consideration to more important processes. Therefore, it would have been obvious to a person of ordinary skill in the art at the time the invention was made to incorporate the round robin method of OS-9 in the device to ensure all processes are executed while maintaining priority and simplicity.

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- 9. Referring to claim 8, Parady has taught for use with a multi-thread execution pipeline loop having a pipeline latency, a method of operating a context switching system, comprising:
 - a. Detecting a device request from a thread executing within said multi-thread execution pipeline loop for access to a device having a fulfillment latency exceeding said pipeline latency (Parady Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3);
 - b. Generating a context switch request for said thread when said thread issues said device request (Parady Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3); and
 - c. Receiving said context switch request (Parady Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3).

executed while maintaining priority and simplicity.

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10. Parady has not explicitly taught storing said thread based thereon in a miss fulfillment first-in-first-out buffer (FIFO) until said device request is fulfilled, said thread sequencing through said entire miss fulfillment FIFO before exiting therefrom. However, Parady has taught using a round robin scheduling method (Parady column 4, lines 9-11). OS-9 has taught round robin scheduling (OS-9 page 5) storing said thread based thereon in a miss fulfillment first-infirst-out buffer (FIFO) until said device request is fulfilled, said thread sequencing through said entire miss fulfillment FIFO before exiting therefrom (OS-9 page 5). In regards to OS-9, the event queue stores processes that are blocked, waiting for some event to complete, like an I/O access. This is similar to a miss fulfillment FIFO, which, by its name and the language in the claim, stores processes that wait for some other event, like an I/O access, to complete. A person of ordinary skill in the art at the time the invention was made would have recognized that a round-robin scheduling method is a fairly simple method that guarantees that all processes will be executed while ensuring that priority between processes is maintained, thereby preventing starvation of a process while giving consideration to more important processes. Therefore, it would have been obvious to a person of ordinary skill in the art at the time the invention was made to incorporate the round robin method of OS-9 in the device to ensure all processes are

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11. Referring to claims 2 and 9, Parady in view of OS-9 has taught wherein said context controller subsystem is further configured to allow a new thread to enter said multi-thread execution pipeline loop after storing said thread in said misfulfillment FIFO (OS-9 page 5). In regards to OS-9, after waiting for the event, e.g. I/O access, to complete, the process is moved to

the active queue, which holds processes active and ready to be executed (OS-9 page 5). From the active queue, the highest priority process runs each tick (OS-9 page 5).

- 12. Referring to claims 3 and 10, Parady in view of OS-9 has taught wherein said context controller subsystem is further configured to allow other threads within said multi-thread execution pipeline loop to continue to execute while said thread is waiting for said device request to be fulfilled (Parady column 2, lines 28-34; column 3, line 57 to column 4, line 18; column 4, line 42-62; and Figure 3).
- 13. Referring to claims 4 and 11, Parady in view of OS-9 has taught wherein said context controller subsystem is further configured to:
 - a. Store said thread in said miss fulfillment FIFO upon reaching an end position of said multi-thread execution pipeline loop (OS-9 page 5), and
 - b. Reinsert said thread into said multi-thread execution pipeline loop at a beginning position after said thread exits said miss fulfillment FIFO (OS-9 page 5).
- 14. Referring to claims 5 and 12, Parady in view of OS-9 has taught wherein said thread is looped back to a beginning stage of said multi-thread execution pipeline loop when said thread reaches an end stage of said multi-thread execution pipeline loop and said thread has not finished processing (OS-9 page 5).
- 15. Referring to claims 6 and 13, Parady in view of OS-9 has taught wherein said context controller subsystem is further configured to sequence said thread through said miss fulfillment FIFO at a rate having a period substantially equivalent to said pipeline latency (OS-9 page 5). In regards to OS-9, the wait queue only holds processes that are blocked and waiting for some event

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to complete, e.g. waiting for an I/O access to complete. This I/O access latency substantially equivalent to the pipeline latency.

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- 16. Referring to claims 7 and 14, Parady in view of OS-9 has taught wherein said device request is a request to access external memory due to a cache miss status (Parady Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3).
- 17. Claims 15-17 and 20-21 are rejected under 35 U.S.C. 103(a) as being unpatentable over Parady, U.S. Patent Number 5,933,627 (herein referred to as Parady) in view of Wilford et al., U.S. Patent Number 5,509,006 (herein referred to as Wilford) and in further view of MicrowareTM's "OS-9: Real-Time Operating System The Complete Software Solution for Your Embedded Application" ©1999 (herein referred to as OS-9).
- 18. Referring to claim 15, Parady has taught a fast pattern processor that receives and processes protocol data units (PDUs), comprising:
 - a. A dynamic random access memory (DRAM) that contains instructions (Parady column 5, lines 19-22; Figure 5; and Figure 6). In regards to Parady, DRAM in a specific type of RAM and Parady shows that RAM is used in his system. Please see Rosenberg's Computers, Information Processing & Telecommunications

 Second Edition for more information of RAM and DRAM.
 - b. A memory cache that caches certain of said instructions from said DRAM (Parady column 5, lines 19-22; Figure 5; and Figure 6); and
 - c. An engine that employs said DRAM and said memory cache to obtain ones of said instructions (Parady Abstract; column 1, lines 29-35 and 46-57; column 2,

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lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3), including:

- i. A multi-thread execution pipeline loop having a pipeline latency (Parady Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3), and
- ii. A context switching system for said multi-thread execution pipeline loop (Parady Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3), having:
 - (1) A context switch requesting subsystem that: detects a device request from a thread executing within said multi-thread execution pipeline loop for access to a device having a fulfillment latency exceeding said pipeline latency (Parady Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3), and
 - (2) Generates a context switch request for said thread (Parady
 Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34;
 column 3, line 57 to column 4, line 18; column 4, lines 42-62;
 Figure 3), and
- iii. A context controller subsystem that receives said context switch request and prevents said thread from executing until said device request is fulfilled (Parady Abstract; column 1, lines 29-35 and 46-57; column 2,

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lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3).

- 19. Parady has not taught a tree engine that parses data within said PDUs. Wilford has taught a tree engine that parses data within said PDUs (Wilford column 1, lines 34-42; column 1, line 65 to column 2, line 19; column 14, lines 14-35; and Figure 5B). A person of ordinary skill in the art at the time the invention was made, and as taught in Wilford, would have recognized that a tree engine that parses data within said PDUs identifies which protocol the data belongs to in order to send the data to the correct destination (Wilford column 1, lines 34-42), thereby ensuring correct data execution. Therefore, it would have been obvious to a person of ordinary skill in the art to incorporate the tree engine of Wilford in the device of Parady to ensure correct data execution.
- 20. In addition, Parady has not explicitly taught
 - a. A miss fulfillment first-in-first-out buffer (FIFO); and
 - b. Based thereon, stores said thread in said miss fulfillment FIFO until said device request is fulfilled, said thread sequencing through said entire miss fulfillment FIFO before exiting therefrom.
- 21. However, Parady has taught using a round robin scheduling method (Parady column 4, lines 9-11). OS-9 has taught round robin scheduling (OS-9 page 5) with a
 - a. A miss fulfillment first-in-first-out buffer (FIFO) (OS-9 page 5); and
 - b. Based thereon, stores said thread in said miss fulfillment FIFO until said device request is fulfilled, said thread sequencing through said entire miss fulfillment FIFO before exiting therefrom (OS-9 page 5).

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22. In regards to OS-9, the event queue stores processes that are blocked, waiting for some event to complete, like an I/O access. This is similar to a miss fulfillment FIFO, which, by its name and the language in the claim, stores processes that wait for some other event, like an I/O access, to complete. A person of ordinary skill in the art at the time the invention was made would have recognized that a round-robin scheduling method is a fairly simple method that guarantees that all processes will be executed while ensuring that priority between processes is maintained, thereby preventing starvation of a process while giving consideration to more important processes. Therefore, it would have been obvious to a person of ordinary skill in the art at the time the invention was made to incorporate the round robin method of OS-9 in the device to ensure all processes are executed while maintaining priority and simplicity.

- 23. Referring to claim 16, Parady in view of Wilford and in further view of OS-9 has taught wherein said context controller subsystem further allows a new thread to enter said multi-thread execution pipeline loop after storing said thread in said FIFO (OS-9 page 5). In regards to OS-9, after waiting for the event, e.g. I/O access, to complete, the process is moved to the active queue, which holds processes active and ready to be executed (OS-9 page 5). From the active queue, the highest priority process runs each tick (OS-9 page 5).
- 24. Referring to claim 17, Parady in view of Wilford and in further view of OS-9 has taught wherein said context controller subsystem further allows other threads within said multi-thread execution pipeline loop to continue to execute while said thread is waiting for said device request to be fulfilled (Parady column 2, lines 28-34; column 3, line 57 to column 4, line 18; column 4, line 42-62; and Figure 3).

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25. Referring to claim 18, Parady in view of Wilford and in further view of OS-9 has taught wherein said context controller subsystem is further configured to:

- d. Store said thread in said miss fulfillment FIFO upon reaching an end position of said multi-thread execution pipeline loop (OS-9 page 5), and
- e. Reinsert said thread into said multi-thread execution pipeline loop at a beginning position after said thread exits said miss fulfillment FIFO (OS-9 page 5).
- 26. Referring to claim 19, Parady in view of Wilford and in further view of OS-9 has taught wherein said thread is looped back to a beginning stage of said multi-thread execution pipeline loop when said thread reaches an end stage of said multi-thread execution pipeline loop and said thread has not finished processing (OS-9 page 5).
- 27. Referring to claim 20, Parady in view of Wilford and in further view of OS-9 has taught wherein said context controller subsystem further sequences said thread through said miss fulfillment FIFO at a rate having a period substantially equivalent to said pipeline latency (OS-9 page 5). In regards to OS-9, the wait queue only holds processes that are blocked and waiting for some event to complete, e.g. waiting for an I/O access to complete. This I/O access latency substantially equivalent to the pipeline latency.
- 28. Referring to claim 21, Parady in view of Wilford and in further view of OS-9 has taught wherein said device request is said DRAM and said device request is a request to access said DRAM due to a cache miss status from said memory cache (Parady Abstract; column 1, lines 29-35 and 46-57; column 2, lines 18-34; column 3, line 57 to column 4, line 18; column 4, lines 42-62; Figure 3).

Response to Arguments

persuasive. Applicants' argue in essence on pages 7-9

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29. Applicant's arguments filed 23 June 2006 have been fully considered but they are not

...unlike the miss fulfillment FIFO of amended Claims 1 and 8, processes that enter the queues of OS-9 do not sequence through the entire queues before exiting therefrom. Instead, the queues of OS-9 are general storage areas that hold a process until a child process terminates, the process is awakened, an event value is received or a semaphore is received...

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30. This has not been found persuasive. Applicants' arguments seem to rely on statements which the Examiner could not locate within OS-9. OS-9 states at most "The highest priority process runs each tick. The lower priority process returns to the active queue." OS-9 also states "All unblocked processes go back on the active queue before they can run again." There is no support in these statements supporting the allegation that the queues are accessed out of turn. In fact, the language actually insinuates that the lower priority processes are removed from the queue, since the "lower priority process returns to the active queue." Also, if OS-9 actually meant for the elements to be removed prior to a full cycle through the entire FIFO, then OS-9 would have called its queues "buffers", "caches", "RAM", etc., since the term "queue" has a very set and well-known definition in the art. As taught in the definition of "queue" the Examiner provided in a previous Office Action, a "queue" is "a first-in first-out data structure used to sequence multiple demands for a resource...Objects are added to the tail of the queue and taken off the head." This means that, in a "queue", elements are added to the tail, i.e. end, of the queue, and taken from the head, i.e. beginning, of the queue, so the elements have to sequence through the entire queue since it is added to the end and removed from the beginning, like a

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standard line at the bank or store and as claimed. To further support this position, the definition

of "first-in first-out", also known as FIFO and "queue" as stated in the definition, was provided

as well (FOLDOC "first-in first-out" ©1999). The definition of FIFO explicitly states "a data

structure or hardware buffer from which items are taken out in the same order they were put in."

This means that the data has to sequence through the entire structure in order for the elements to

be taken out in the same order they were put in.

Conclusion

31. Any inquiry concerning this communication or earlier communications from the

examiner should be directed to Aimee J. Li whose telephone number is (571) 272-4169. The

examiner can normally be reached on M-T 7:30am-5:00pm.

32. If attempts to reach the examiner by telephone are unsuccessful, the examiner's

supervisor, Eddie Chan can be reached on (571) 272-4162. The fax phone number for the

organization where this application or proceeding is assigned is 703-872-9306.

33. Information regarding the status of an application may be obtained from the Patent

Application Information Retrieval (PAIR) system. Status information for published applications

may be obtained from either Private PAIR or Public PAIR. Status information for unpublished

applications is available through Private PAIR only. For more information about the PAIR

system, see http://pair-direct.uspto.gov. Should you have questions on access to the Private PAIR

system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free).

AJL Aimee J. Li 31 July 2006 Solei CHAN

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